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## ALMOST $n$ -ARY AND ALMOST $n$ -ARITIZABLE THEORIES

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**ABSTRACT.** We study possibilities for almost  $n$ -ary and  $n$ -aritzable theories. Their dynamics in general case, for  $\omega$ -categorical theories, and with respect to operations for theories are described.

**Keywords:** elementary theory, almost  $n$ -ary theory, almost  $n$ -aritzable theory.

We continue to study arities of theories and of their expansions [1]. In the present paper we introduce natural notions of almost  $n$ -ary and almost  $n$ -aritzable theories, and describe their dynamics both in general case, for  $\omega$ -categorical theories, and with respect to operations for theories.

### 1. PRELIMINARIES

Recall a series of notions related to arities and aritzabilities of theories.

**Definition** [2]. A theory  $T$  is said to be  $\Delta$ -based, where  $\Delta$  is some set of formulae without parameters, if any formula of  $T$  is equivalent in  $T$  to a Boolean combination of formulae in  $\Delta$ .

For  $\Delta$ -based theories  $T$ , it is also said that  $T$  has *quantifier elimination* or *quantifier reduction* up to  $\Delta$ .

**Definition** [2, 3]. Let  $\Delta$  be a set of formulae of a theory  $T$ , and  $p(\bar{x})$  a type of  $T$  lying in  $S(T)$ . The type  $p(\bar{x})$  is said to be  $\Delta$ -based if  $p(\bar{x})$  is isolated by a set of formulas  $\varphi^\delta \in p$ , where  $\varphi \in \Delta$ ,  $\delta \in \{0, 1\}$ .

The following lemma, being a corollary of Compactness Theorem, noticed in [2].

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**Lemma 1.** *A theory  $T$  is  $\Delta$ -based if and only if, for any tuple  $\bar{a}$  of any (some) weakly saturated model of  $T$ , the type  $\text{tp}(\bar{a})$  is  $\Delta$ -based.*

**Definition [1].** An elementary theory  $T$  is called *unary*, or *1-ary*, if any  $T$ -formula  $\varphi(\bar{x})$  is  $T$ -equivalent to a Boolean combination of  $T$ -formulas, each of which is of one free variable, and of formulas of form  $x \approx y$ .

For a natural number  $n \geq 1$ , a formula  $\varphi(\bar{x})$  of a theory  $T$  is called  *$n$ -ary*, or an  *$n$ -formula*, if  $\varphi(\bar{x})$  is  $T$ -equivalent to a Boolean combination of  $T$ -formulas, each of which is of  $n$  free variables.

For a natural number  $n \geq 2$ , an elementary theory  $T$  is called  *$n$ -ary*, or an  *$n$ -theory*, if any  $T$ -formula  $\varphi(\bar{x})$  is  $n$ -ary.

A theory  $T$  is called *binary* if  $T$  is 2-ary, it is called *ternary* if  $T$  is 3-ary, etc.

We will admit the case  $n = 0$  for  $n$ -formulae  $\varphi(\bar{x})$ . In such a case  $\varphi(\bar{x})$  is just  $T$ -equivalent to a sentence  $\forall \bar{x} \varphi(\bar{x})$ .

If  $T$  is a theory such that  $T$  is  $n$ -ary and not  $(n - 1)$ -ary then the value  $n$  is called the *arity* of  $T$  and it is denoted by  $\text{ar}(T)$ . If  $T$  does not have any arity we put  $\text{ar}(T) = \infty$ .

Similarly, for a formula  $\varphi$  of a theory  $T$  we denote by  $\text{ar}_T(\varphi)$  the natural value  $n$  if  $\varphi$  is  $n$ -ary and not  $(n - 1)$ -ary. If  $\varphi$  does not have any arity we put  $\text{ar}_T(\varphi) = \infty$ . If a theory  $T$  is fixed we write  $\text{ar}(\varphi)$  instead of  $\text{ar}_T(\varphi)$ .

The following example illustrates the notions above, and it will be used below.

**Example 1.** Recall [4, 5, 6] that a *circular*, or *cyclic* order relation is described by a ternary relation  $K_3$  satisfying the following conditions:

- (co1)  $\forall x \forall y \forall z (K_3(x, y, z) \rightarrow K_3(y, z, x))$ ;
- (co2)  $\forall x \forall y \forall z (K_3(x, y, z) \wedge K_3(y, x, z) \leftrightarrow x = y \vee y = z \vee z = x)$ ;
- (co3)  $\forall x \forall y \forall z (K_3(x, y, z) \rightarrow \forall t [K_3(x, y, t) \vee K_3(t, y, z)])$ ;
- (co4)  $\forall x \forall y \forall z (K_3(x, y, z) \vee K_3(y, x, z))$ .

Clearly,  $\text{ar}(K_3(x, y, z)) = 3$  if the relation has at least three element domain. Hence, theories with infinite circular order relations are at least 3-ary.

The following generalization of circular order produces a  *$n$ -ball*, or  *$n$ -spherical*, or  *$n$ -circular* order relation, for  $n \geq 4$ , which is described by a  $n$ -ary relation  $K_n$  satisfying the following conditions:

- (nbo1)  $\forall x_1, \dots, x_n (K_n(x_1, x_2, \dots, x_n) \rightarrow K_n(x_2, \dots, x_n, x_1))$ ;

- (nbo2)  $\forall x_1, \dots, x_n \left( K_n(x_1, \dots, x_i, x_{i+1}, \dots, x_n) \wedge \right.$   
 $\left. \wedge K_n(x_1, \dots, x_{i+1}, x_i, \dots, x_n) \leftrightarrow \bigvee_{i=1}^{n-1} x_i = x_{i+1} \right)$ ;

- (nbo3)  $\forall x_1, \dots, x_n (K_n(x_1, \dots, x_n) \rightarrow \forall t [K_n(x_1, \dots, x_{n-1}, t) \vee K_n(t, x_2, \dots, x_n)])$ ;

- (nbo4)  $\forall x_1, \dots, x_n (K_n(x_1, \dots, x_i, x_{i+1}, \dots, x_n) \vee K_n(x_1, \dots, x_{i+1}, x_i, \dots, x_n))$ ,  $i < n$ .

Clearly,  $\text{ar}(K_n(x_1, \dots, x_n)) = n$  if the relation has at least  $n$ -element domain. Thus, theories with infinite  $n$ -ball order relations are at least  $n$ -ary.

**Definition [1].** A  $T$ -formula  $\varphi(\bar{x})$  is called  *$n$ -expansible*, or  *$n$ -arizable*, or  *$n$ -aritzable*, if  $T$  has an expansion  $T'$  such that  $\varphi(\bar{x})$  is  $T'$ -equivalent to a Boolean combination of  $T'$ -formulas with  $n$  free variables.

A theory  $T$  is called  $n$ -*expansible*, or  $n$ -*arizable*, or  $n$ -*aritizable*, if there is an  $n$ -ary expansion  $T'$  of  $T$ .

A theory  $T$  is called *arizable* or *aritizable*, if  $T$  is  $n$ -aritizable for some  $n$ .

A 1-aritizable theory is called *unary-able*, or *unary-tizable*. A 2-aritizable theory is called *binary-tizable* or *binarizable*, a 3-aritizable theory is called *ternary-tizable* or *ternarizable*, etc.

**Definition.** [8] The *disjoint union*  $\bigsqcup_{n \in \omega} \mathcal{M}_n$  of pairwise disjoint structures  $\mathcal{M}_n$  for pairwise disjoint predicate languages  $\Sigma_n$ ,  $n \in \omega$ , is the structure of language  $\bigcup_{n \in \omega} \Sigma_n \cup \{P_n^{(1)} \mid n \in \omega\}$  with the universe  $\bigsqcup_{n \in \omega} M_n$ ,  $P_n = M_n$ , and interpretations of predicate symbols in  $\Sigma_n$  coinciding with their interpretations in  $\mathcal{M}_n$ ,  $n \in \omega$ . The *disjoint union of theories*  $T_n$  for pairwise disjoint languages  $\Sigma_n$  accordingly,  $n \in \omega$ , is the theory

$$\bigsqcup_{n \in \omega} T_n \Leftrightarrow \text{Th} \left( \bigsqcup_{n \in \omega} \mathcal{M}_n \right),$$

where  $\mathcal{M}_n \models T_n$ ,  $n \in \omega$ . Taking empty sets instead of some structures  $\mathcal{M}_k$  we obtain disjoint unions of finitely many structures and theories. In particular, we have the disjoint unions  $\mathcal{M}_0 \sqcup \dots \sqcup \mathcal{M}_n$  and their theories  $T_0 \sqcup \dots \sqcup T_n$ .

**Theorem 1.** [1]. 1. For any theories  $T_m$ ,  $m \in \omega$ , and their disjoint union  $\bigsqcup_{m \in \omega} T_m$ , all  $T_m$  are  $n$ -theories iff  $\bigsqcup_{m \in \omega} T_m$  is an  $n$ -theory, moreover,

$$\text{ar} \left( \bigsqcup_{m \in \omega} T_m \right) = \max \{ \text{ar}(T_m) \mid m \in \omega \}.$$

2. For any theories  $T_m$ ,  $m \in \omega$ , and their disjoint union  $\bigsqcup_{m \in \omega} T_m$ , all  $T_m$  are  $n$ -aritizable iff  $\bigsqcup_{m \in \omega} T_m$  is  $n$ -aritizable.

**Definition** [7]. Let  $\mathcal{M}$  and  $\mathcal{N}$  be structures of relational languages  $\Sigma_{\mathcal{M}}$  and  $\Sigma_{\mathcal{N}}$  respectively. We define the *composition*  $\mathcal{M}[\mathcal{N}]$  of  $\mathcal{M}$  and  $\mathcal{N}$  satisfying the following conditions:

- 1)  $\Sigma_{\mathcal{M}[\mathcal{N}]} = \Sigma_{\mathcal{M}} \cup \Sigma_{\mathcal{N}}$ ;
- 2)  $M[\mathcal{N}] = M \times N$ , where  $M[\mathcal{N}]$ ,  $M$ ,  $N$  are universes of  $\mathcal{M}[\mathcal{N}]$ ,  $\mathcal{M}$ , and  $\mathcal{N}$  respectively;
- 3) if  $R \in \Sigma_{\mathcal{M}} \setminus \Sigma_{\mathcal{N}}$ ,  $\mu(R) = n$ , then  $((a_1, b_1), \dots, (a_n, b_n)) \in R_{\mathcal{M}[\mathcal{N}]}$  if and only if  $(a_1, \dots, a_n) \in R_{\mathcal{M}}$ ;
- 4) if  $R \in \Sigma_{\mathcal{N}} \setminus \Sigma_{\mathcal{M}}$ ,  $\mu(R) = n$ , then  $((a_1, b_1), \dots, (a_n, b_n)) \in R_{\mathcal{M}[\mathcal{N}]}$  if and only if  $a_1 = \dots = a_n$  and  $(b_1, \dots, b_n) \in R_{\mathcal{N}}$ ;
- 5) if  $R \in \Sigma_{\mathcal{M}} \cap \Sigma_{\mathcal{N}}$ ,  $\mu(R) = n$ , then  $((a_1, b_1), \dots, (a_n, b_n)) \in R_{\mathcal{M}[\mathcal{N}]}$  if and only if  $(a_1, \dots, a_n) \in R_{\mathcal{M}}$ , or  $a_1 = \dots = a_n$  and  $(b_1, \dots, b_n) \in R_{\mathcal{N}}$ .

The composition  $\mathcal{M}[\mathcal{N}]$  is called *e-definable*, or *equ-definable*, if  $\mathcal{M}[\mathcal{N}]$  has an  $\emptyset$ -definable equivalence relation  $E$  whose  $E$ -classes are universes of the copies of  $\mathcal{N}$  forming  $\mathcal{M}[\mathcal{N}]$ . If the equivalence relation  $E$  is fixed, the *e-definable composition* is called *E-definable*.

Using a nice basedness of  $E$ -definable compositions  $T_1[T_2]$  (see [7]) till the formulas of form  $E(x, y)$  and generating formulas for  $T_1$  and  $T_2$  we have the following:

**Theorem 2.** [1]. 1. For any theories  $T_1$  and  $T_2$  and their  $E$ -definable composition  $T_1[T_2]$ ,  $T_1$  and  $T_2$  are  $n$ -theories, for  $n \geq 2$ , iff  $T_1[T_2]$  is an  $n$ -theory, moreover,  $\text{ar}(T_1[T_2]) = \max\{\text{ar}(T_1), \text{ar}(T_2)\}$ , if models of  $T_1$  and of  $T_2$  have at least two elements, and  $\text{ar}(T_1[T_2]) = \max\{\text{ar}(T_1), \text{ar}(T_2), 2\}$ , if a model of  $T_1$  or  $T_2$  is a singleton.

2. For any theories  $T_1$  and  $T_2$  and their  $E$ -definable composition  $T_1[T_2]$ ,  $T_1$  and  $T_2$  are  $n$ -aritzable iff  $T_1[T_2]$  is  $n$ -aritzable.

## 2. ALMOST $n$ -ARY AND $n$ -ARITIZABLE THEORIES, THEIR DYNAMICS

**Definition.** (Cf. [5, 6]) A theory  $T$  is called *almost  $n$ -ary* if there are finitely many formulae  $\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})$  such that each  $T$ -formula is  $T$ -equivalent to a Boolean combination of  $n$ -formulae and formulae obtained by substitutions of free variables in  $\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})$ .

In such a case we say that the formulae  $\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})$  witness that  $T$  is almost  $n$ -ary.

Almost 1-ary theories are called *almost unary*, almost 2-ary theories are called *almost binary*, almost 3-ary theories are called *almost ternary*, etc.

A theory  $T$  is called *almost  $n$ -aritzable* if some expansion  $T'$  of  $T$  is almost  $n$ -ary.

Almost 1-aritzable theories are called *almost unary-tizable*, almost 2-aritzable theories are called *almost binarizable*, almost 3-aritzable theories are called *almost ternarizable*, etc.

The following properties are obvious.

1. Any  $n$ -ary (respectively,  $n$ -aritzable) theory is almost  $n$ -ary (almost  $n$ -aritzable).
2. Any almost  $n$ -ary (respectively,  $n$ -aritzable) theory is almost  $k$ -ary (almost  $k$ -aritzable) for any  $k \geq n$ .
3. Any theory of a finite structure is almost unary.

Families of weakly circularly minimal structures produce examples of almost binary theories which are not binary [5, 6]. Similarly natural generalizations of weakly circularly minimal structures till  $n$ -circular orders give examples of almost  $(n-1)$ -ary theories  $T_n$  with  $\text{ar}(T_n) = n$ ,  $n \geq 4$ .

Assuming that the witnessing set  $\{\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})\}$  is minimal for the almost  $n$ -ary theory  $T$  we have either  $m = 0$  or  $l(\bar{x}) > n$ .

Thus we have two minimal characteristics witnessing the almost  $n$ -arity of  $T$ :  $m$  and  $l(\bar{x})$ . The pair  $(m, l(\bar{x}))$  is called the *degree* of the almost  $n$ -arity of  $T$ , or the *aar-degree* of  $T$ , denoted by  $\text{deg}_{\text{aar}}(T)$ . Here we assume that  $n$  is minimal with almost  $n$ -arity of  $T$ , this  $n$  is denoted by  $\text{aar}(T)$ . Clearly,  $\text{aar}(T) \leq \text{ar}(T)$ , and if  $m = 0$ , i.e.,  $n = \text{ar}(T) = \text{aar}(T)$  then it is supposed that  $l(\bar{x}) = 0$ , too.

We have  $\text{aar}(T) \in \omega$  if and only if  $\text{ar}(T) \in \omega$ . So if  $\text{ar}(T) = \infty$  then it is natural to put  $\text{aar}(T) = \infty$ .

Besides,  $n = \text{ar}(T) = \text{aar}(T) \in \omega$  means that the set  $\Delta_n(T)$  of  $T$ -formulae with  $n$  free variables allows to express all  $\emptyset$ -definable sets for  $T$  by Boolean combinations and taking any set  $\Delta_k(T)$  of  $T$ -formulae with  $k < n$  free variables all  $\emptyset$ -definable sets for  $T$  can be expressed by Boolean combinations of formulae in  $\Delta_k$  and substitutions of infinitely many formulae  $\varphi(\bar{x})$  only, where  $l(\bar{x}) = n$ .

By the definition if  $\text{aar}(T) = n$  then

$$(1) \quad \text{deg}_{\text{aar}}(T) \in \{(0, 0)\} \cup \{(m, r) \mid m \in \omega \setminus \{0\}, r \in \omega, r > n\}.$$

The described pairs in the relation (1) are called *admissible*.

**Theorem 3.** *For any  $m, n \in \omega \setminus \{0\}$  with  $m \leq n$  there is a theory  $T_{mn}$  with  $\text{aar}(T_{mn}) = m$  and  $\text{ar}(T_{mn}) = n$ .*

Proof. We use disjoint unions of *dense  $n$ -spherically ordered theories*  $T_n$  with infinite orders, i.e., theories, generated by  $n$ -spherical orders  $K_n(x_1, x_2, x_3, \dots, x_n)$  satisfying the axioms

$$\begin{aligned} \forall x_1, x_2, \dots, x_n \left( K_n(x_1, x_2, x_3, \dots, x_n) \wedge \bigwedge_{i \neq j} \neg x_i \approx x_j \rightarrow \right. \\ \left. \rightarrow \exists y \left( \bigwedge_{i \leq n} \neg x_i \approx y \wedge K_n(x_1, y, x_3, \dots, x_n) \right) \right). \end{aligned}$$

For  $n = 2$  we take the theory  $T_2$  of dense linear order  $K_2(x_1, x_2)$  without endpoints, having  $\text{ar}(T_2) = 2$ , and for  $n = 1$  — the theory  $T_1$  of the empty languages, having  $\text{ar}(T_1) = 1$ .

Similarly to dense linear orders and dense circular orders, dense  $n$ -spherical orders produce quantifier eliminations with  $\text{ar}(T_n) = n$ ,  $n \in \omega \setminus \{0\}$ .

Using Theorem 1 we obtain  $m = \text{ar}(T) = \text{aar}(T)$  taking  $T_{mm} = \bigsqcup_{r \in \omega} T_m^r$ ,

where  $T_m^r$  are copies of  $T_m$  in disjoint languages  $\{K_m^r\}$ ,  $r \in \omega$ . Indeed, disjoint predicates  $K_m^r$  producing  $\text{ar}(T_m) = m$  witness that  $\text{ar}(T_{mm}) = m$ . Finitely many these predicates can not define all definable sets for  $T_{mm}$  since there are infinitely many of them. Thus,  $\text{aar}(T_{mm}) = m$ , too.

Now for any  $m < n$  we form  $T_{mn} = T_n \sqcup \bigsqcup_{r \in \omega} T_m^r$ . By  $T_n$  in  $T_{mn}$  and  $m < n$  we have  $\text{ar}(T_{mn}) = n$ . And  $\text{aar}(T_{mn}) = m$  since there are infinitely many disjoint predicates of arity  $m$ .

The following theorem shows that all admissible pairs are realized.

**Theorem 4.** *For any admissible pair  $(m, r)$  and  $n \in \omega \setminus \{0, 1\}$  there is a theory  $T$  with  $\text{aar}(T) = n$  and  $\text{deg}_{\text{ar}}(T) = (m, r)$ .*

Proof. The admissible pair  $(0, 0)$  with  $\text{ar}(T) = n$  is realized by Theorem 3. Now for an admissible pair  $(m, r) \neq (0, 0)$  we can take a disjoint union  $T$  of countably many dense  $n$ -spherically ordered theories and of  $m$  dense  $r$ -spherically ordered theories. Using arguments for Theorem 3 we obtain  $\text{aar}(T) = n$  and  $\text{deg}_{\text{ar}}(T) = (m, r)$ .  $\square$

**Proposition 1.** *Any almost  $n$ -ary theory  $T$  is  $k$ -ary for some  $k \geq n$ .*

Proof. Let  $T$  be an almost  $n$ -ary theory witnessed by the formulae  $\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})$ . Then taking the set  $\Delta_k$  of all formulae with  $k = \max\{n, l(\bar{x})\}$  we observe, using  $\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})$ , that  $T$  is  $\Delta_k$ -based, i.e.,  $T$  is  $k$ -ary.

**Corollary 1.** *Any theory  $T$  is  $n$ -ary for some  $n$  iff  $T$  is almost  $m$ -ary for some  $m$ .*

**Corollary 2.** *Any theory  $T$  is  $n$ -aritzable for some  $n$  iff  $T$  is almost  $m$ -aritzable for some  $m$ .*

3.  $\omega$ -CATEGORICAL ALMOST  $n$ -ARY AND  $n$ -ARITIZABLE THEORIES

**Proposition 2.** *If  $T$  is an almost  $n$ -ary  $\omega$ -categorical theory, for some  $n$ , then  $T$  is almost  $k$ -ary for any  $k \in \omega \setminus \{0\}$ , i.e.,  $\text{aar}(T) = 1$ .*

Proof. If  $k \geq n$  then  $T$  is almost  $m$ -ary, as noticed above. If  $k < n$  then we collect in a set  $Z$  all formulae  $\varphi_1(\bar{x}), \dots, \varphi_m(\bar{x})$  witnessing the almost  $n$ -arity of  $T$  and, by Ryll-Nardzewski Theorem, all non-equivalent formulae  $\psi_1(\bar{y}), \dots, \psi_r(\bar{y})$  with  $l(\bar{y}) = n$ . Clearly, the set  $Z$  witnesses that  $T$  is almost  $k$ -ary. Taking  $k = 1$  we obtain  $\text{aar}(T) = 1$ .

**Corollary 3.** *For any  $\omega$ -categorical theory  $T$  either  $\text{aar}(T) = 1$  with  $\text{ar}(T) \in \omega$ , or  $\text{aar}(T) = \infty$  with  $\text{ar}(T) = \infty$ .*

The following example illustrates Corollary 3.

**Example 2.** Taking a dense linear order  $K_2$  without endpoint we can step-by-step extend it to a chain of dense  $n$ -spherical orders  $K_n$ ,  $n \geq 2$ , in the following way.

We put  $(a, b, c) \in K_3$  if  $(a, b) \in K_2$  and  $(b, c) \in K_2$ , or  $(b, c) \in K_2$  and  $(c, a) \in K_2$ , or  $(c, a) \in K_2$  and  $(a, b) \in K_2$ . If  $K_n$ ,  $n \geq 3$ , is defined then we put  $(a_1, \dots, a_{n+1}) \in K_{n+1}$  if  $(a_1, \dots, a_n) \in K_n$  and  $(a_2, \dots, a_{n+1}) \in K_n$ , or  $a_2, \dots, a_{n+1} \in K_n$  and  $(a_3, \dots, a_{n+1}, a_1) \in K_n$ , or  $(a_3, \dots, a_{n+1}, a_1) \in K_n$  and  $(a_4, \dots, a_{n+1}, a_1, a_2) \in K_n$ . The obtained structure  $\mathcal{M}_\infty$  in the language  $\Sigma_\infty = \{K_n \mid n \in \omega \setminus \{0, 1\}\}$  has an  $\omega$ -categorical theory  $T_\infty$  with  $\text{aar}(T_\infty) = \text{ar}(T_\infty) = \infty$ , since each new  $K_n$  increases the arity. The same characteristics have restrictions of  $T_\infty$  to any infinite sublanguages.

At the same time each restriction  $\mathcal{M}$  of  $\mathcal{M}_\infty$  to a finite nonempty sublanguage  $\{K_{n_1}, \dots, K_{n_m}\}$  produces a theory  $T$  with  $\text{aar}(T) = 1$  with

$$\text{ar}(T) = \max\{n_1, \dots, n_m\}.$$

By Corollary 3 and the definition of almost aritizability we immediately have:

**Corollary 4.** *Any restriction  $T$  of  $n$ -ary  $\omega$ -categorical theory  $T'$  is almost unary-aritizable.*

4. OPERATIONS FOR ALMOST  $n$ -ARY AND  $n$ -ARITIZABLE THEORIES

In this section we consider links for arities of theories with respect to disjoint unions of theories and  $E$ -definable compositions of theories.

**Theorem 5.** 1. *For any theories  $T_1, T_2$  and their disjoint union  $T_1 \sqcup T_2$ , both  $T_1$  and  $T_2$  are almost  $n$ -ary iff  $T_1 \sqcup T_2$  is almost  $n$ -ary, moreover,  $\text{aar}(T_1 \sqcup T_2) = \max\{\text{aar}(T_1), \text{aar}(T_2)\}$ .*

2. *For any theories  $T_1, T_2$  and their disjoint union  $T_1 \sqcup T_2$ , both  $T_1$  and  $T_2$  are almost  $n$ -aritizable iff  $T_1 \sqcup T_2$  is almost  $n$ -aritizable.*

Proof. 1. Let  $\Phi_1$  and  $\Phi_2$  be finite sets of formulas witnessing that  $T_1$  and  $T_2$  are almost  $n$ -ary, respectively. Using the definition of disjoint union and Theorem 1 we obtain that the finite set  $\Phi_1 \cup \Phi_2$  witnesses that  $T_1 \sqcup T_2$  is almost  $n$ -ary. Conversely, if a finite set  $\Phi$  of formulas witnesses that  $T_1 \sqcup T_2$  is almost  $n$ -ary then  $\Phi$  witnesses that both  $T_1$  and  $T_2$  are almost  $n$ -ary.

The equality  $\text{aar}(T_1 \sqcup T_2) = \max\{\text{aar}(T_1), \text{aar}(T_2)\}$  follows from the definition of disjoint union since if  $\text{aar}(T_1 \sqcup T_2) = k$  then the maximal value of  $\text{aar}(T_1)$  and  $\text{aar}(T_2)$  is responsible for this equality.

Item 2 follows from Item 1 since expansions of  $T_1$  and  $T_2$  correspond expansions of  $T_1 \sqcup T_2$ : some expansions of  $T_1'$  and  $T_2'$  of  $T_1$  and  $T_2$ , respectively, are almost  $n$ -ary iff  $T_1' \sqcup T_2'$  produces an almost  $n$ -ary expansion of  $T_1 \sqcup T_2$ .

Using induction we obtain the following:

**Corollary 5.** 1. For any theories  $T_1, T_2, \dots, T_m$  and their disjoint union  $\bigsqcup_{i=1}^m T_i$ , all  $T_1, T_2, \dots, T_m$  are almost  $n$ -ary iff  $\bigsqcup_{i=1}^m T_i$  is almost  $n$ -ary, moreover,

$$\text{aar} \left( \bigsqcup_{i=1}^m T_i \right) = \max\{\text{aar}(T_i) \mid i \leq m\}.$$

2. For any theories  $T_1, T_2, \dots, T_m$  and their disjoint union  $\bigsqcup_{i=1}^m T_i$ , all  $T_1, T_2, \dots, T_m$  are almost  $n$ -aritzable iff  $\bigsqcup_{i=1}^m T_i$  is almost  $n$ -aritzable.

**Remark 1.** Both almost  $n$ -arity and almost  $n$ -aritzability can fail taking disjoint unions of infinitely many theories  $T_i$ ,  $i \in I$ . Indeed, each theory  $T_i$  can have its own finite set  $\Phi_i$  of formulae witnessing the almost  $n$ -arity/ $n$ -aritzability, say in disjoint languages, whereas finite unions  $\bigcup \Phi_i$  can not witness the almost  $n$ -arity/ $n$ -aritzability for  $\bigsqcup_{i \in I} T_i$ .

Generalizing Theorem 2 we obtain:

**Theorem 6.** 1. For any theories  $T_1$  and  $T_2$  and their  $E$ -definable composition  $T_1[T_2]$ ,  $T_1$  and  $T_2$  are almost  $n$ -ary iff  $T_1[T_2]$  is almost  $n$ -ary, moreover,  $\text{aar}(T_1[T_2]) = \max\{\text{aar}(T_1), \text{aar}(T_2), 2\}$ , if models of  $T_1$  and of  $T_2$  have at least two elements, and  $\text{aar}(T_1[T_2]) = \max\{\text{aar}(T_1), \text{aar}(T_2)\}$ , if a model of  $T_1$  or  $T_2$  is a singleton.

2. For any theories  $T_1$  and  $T_2$  and their  $E$ -definable composition  $T_1[T_2]$ ,  $T_1$  and  $T_2$  are almost  $n$ -aritzable iff  $T_1[T_2]$  is almost  $n$ -aritzable.

Proof. 1. Let  $T_i$  be  $\Delta_i$ -based for  $i = 1, 2$ . Since  $T_1[T_2]$  is  $E$ -definable it is  $\Delta$ -based, where  $\Delta$  consists of formulae in  $\Delta_1 \cup \Delta_2$  and  $E(x, y)$  [7]. Now assuming that  $T_1$  and  $T_2$  are almost  $n$ -ary we can choose  $\Delta_i$  consisting of  $n$ -formulae and finitely many formulae forming  $\Phi_i$ ,  $i = 1, 2$ . Hence  $T_1[T_2]$  is almost  $n$ -ary.

Conversely, if  $T_1[T_2]$  is almost  $n$ -ary and it is witnessed by a set  $\Phi$  of formulae then  $\Phi$  witnesses that  $T_1$  and  $T_2$  are almost  $n$ -ary.

If models of  $T_1$  and of  $T_2$  have at least two elements then  $T_1[T_2]$  is at least binary that witnessed by the formula  $E(x, y)$ . Thus since  $T_1[T_2]$  is  $\Delta$ -based we have  $\text{aar}(T_1[T_2]) = \max\{\text{aar}(T_1), \text{aar}(T_2), 2\}$ . If  $T_1$  or  $T_2$  is a theory of singleton then  $T_1[T_2]$  is  $(\Delta_1 \cup \Delta_2)$ -based implying  $\text{aar}(T_1[T_2]) = \max\{\text{aar}(T_1), \text{aar}(T_2)\}$ .

Item 2 follows from Item 1 repeating the arguments for Item 2 of Theorem 5.

Theorem 6 immediately implies

**Corollary 6.** Any composition of finitely many almost  $n$ -ary (almost  $n$ -aritzable) theories, for  $n \geq 2$ , is again many almost  $n$ -ary (almost  $n$ -aritzable).

## 5. CONCLUSION

We considered possibilities for almost arities and almost aritizabilities of theories and their dynamics both in general case, for  $\omega$ -categorical theories, and with respect to operations for theories. It would be interesting to describe values of almost arities and almost aritizabilities for natural classes of theories.

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