

Review of "Bounds for the success probability in the Odds Theorem associated with a given optimal stopping time"

by A. M. Kabaeva, A.V. Logachov and A.A. Yambartsev

The paper under review establishes upper and lower bounds on the success probability in the Bruss Odds Theorem. This theorem solves the following sequential stopping problem. A sequence of independent Bernoulli random variables  $X_1, \dots, X_n$  with respective parameters  $p_1, \dots, p_n$  is observed at time instances  $t = 1, \dots, n$ . The goal is to stop "at the last success" with maximal probability, i.e., to develop a stopping time  $\tau$  such that the probability  $P(X_\tau = 1, X_{\tau+1} = \dots = X_n = 0)$  will be maximal. The Bruss Odds Theorem proves the the optimal stopping rule in this problem is to stop the first time  $s$  such that the sum of odds  $R_s := \sum_{t=s}^n r_t = \sum_{t=s}^n (p_t/q_t)$  is greater than 1; here  $q_t = 1 - p_t$ . Formally,  $s := \max\{1, \max\{1 \leq t \leq n : R_t \geq 1\}\}$ . The corresponding optimal value is

$$V_n = \sum_{t=s}^n r_t \prod_{t=s}^n q_t = R_s \prod_{t=s}^n q_t.$$

The paper under review derives upper and lower bounds on  $V_n$  in terms of  $R_s$ ,  $n$  and  $s$ .

### Remarks

**1.** It is unclear what is the motivation for deriving bounds on the optimal value in terms of  $R_s$ ,  $n$  and  $s$  when the optimal value  $V_n$  is given by a simple explicit formula. The authors write that the success probabilities  $p_1, \dots, p_n$  might be unknown, and then the bounds can be useful. However, in this situation it is unclear how the optimal stopping time  $s$  can be known. In contrast, Bruss (2003) derives a uniform lower bound on  $V_n$ : independently of the problem parameters,  $V_n > 1/e$  whenever  $R_s \geq 1$ .

**2.** The derivation of upper and lower bounds on the optimal value is much more meaningful and interesting in stopping problems in which there are no explicit formulas for  $V_n$  exists.

**3.** It is quite easy to produce tight bounds on  $V_n$  by manipulating the formula for  $V_n$ . Such derivations follow elementary calculations and do not require any probabilistic arguments. The inequality between the arithmetic and geometric means immediately leads to the lower bound

$$V_n = R_s \prod_{t=s}^n q_t = \frac{R_s}{\prod_{t=s}^n (1 + r_t)} \geq \frac{R_s}{(1 + \frac{1}{n-s+1} R_s)^{n-s+1}}.$$

This bound is stated as inequality (2.10) in the proof of item 2 of Theorem 3, it holds for all values of  $R_s$ , not only for the range  $1 \leq R_s \leq 1 + 1/(n - s)$ . This bound is tighter than the one given in the item 2 of Theorem 3. Is there any reason to state a weaker result in the theorem? Also, how this lower bound compares to the one stated in part 3? Anyway, the optimal value is known to satisfy  $V_n \geq R_s e^{-R_s}$  [see Bruss (2000, 2003)], so the differences between the bounds are insignificant.

4. It seems that the authors are not aware of more recent results on the odds theorem and its variants; see, e.g., the following papers and references therein:

Bruss, T. and Louchard, G. (2009). The odds algorithm based on sequential updating and its performance. *Adv. Appl. Probab.* **41**, 131-153.

Dendievel, R. (2013). New developments of the odds-theorem. *Math. Scientist* **38**, 110-122.

Bruss, T. (2019). Odds-theorem and monotonicity. *Mathematica Applicanda* **47**, 25-43.

Goldenshluger, A., Malinovsky, Ya., and Zeevi, A. (2020). A unified approach for solving sequential selection problems. *Probab. Surveys* **17**, 214-256.